## IN THE CLAIMS

Please amend the claims as follows:

1. (Currently Amended) A program storage device, readable by a machine, tangibly embodying programming instructions to perform method steps for constructing a <u>set of types occurring within a method call graph as a representation of a program, the programming instructions comprising:</u>

selecting a program P for constructing a call graph representation thereof; wherein the program P contains zero or more fields  $F_F$  and one or more at least two methods  $M_M$ ;

wherein each method M<sub>1</sub> in M<sub>M</sub> has a single body B;

- wherein for each method M<sub>2</sub> in M<sub>M</sub>, the call graph representation includes a corresponding node;
- wherein the call graph representation includes zero or more edges corresponding to connections between two or more of nodes:

identifying constructing for each method M in  $M_M$ , a set of zero or more- types  $S_M$  of objects which occur therein;

identifying constructing for each field F in  $F_F$ , a set of zero or more types  $S_F$  of objects stored therein;

identifying one or more allocation sites inside the body B of each of method M; determining a set of directly called methods M' that are directly called within inside the body B of each method M and propagating types from the set of types  $S_{M'}$  to the set of types  $S_{M'}$  and from of the set of types  $S_{M}$  to the set of types  $S_{M'}$ ; and

determining a set of virtually called methods M" that are virtually called withininside the body B of each method M and propagating types from the set of types  $S_{\underline{M}}$  to the set of types  $S_{\underline{M}}$  and from the set of types  $S_{\underline{M}}$  to the set of types  $S_{\underline{M}}$ ;

determining a set of fields F that are

read in the body B of each method M, and propagating types from the set

 $\frac{\text{of types }S_F \text{ to the set of types }S_M; \text{ and}}{\text{written in the body B of each method M, and propagating types from the}}$  set of types  $S_M$  to the set of types  $S_F$ .

2. (Currently Amended) The program storage device according to claim 1, further comprising the programming instructions of:

determining the set of all types T that are allocated in the body of method M, and adding each element of the set of all types T to the set of types  $S_M$  for each allocation of type T that occurs in the method M.

3. (Currently Amended) The program storage device according to claim 2, further comprising the programming instructions of:

for each direct call to <u>a</u> method<del>e</del> M' in the body B of the method M performing the steps of:

adding any type that occurs in the <u>set of types</u>  $S_M$  and that is a subtype of the type of a parameter of the methods M' to <u>the set of types</u>  $S_{M'}$ ; and

adding any type that occurs in the <u>set of</u> types  $S_{M'}$  and that is a subtype of a return type of the methods M' to the <u>set of</u> types  $S_{M}$ .

4. (Currently Amended) The program storage device according to claim 3, further comprising the programming instructions of:

for each virtual call to the methods M' in the body B of the method M:

using the <u>set of types</u>  $S_M$ , determine each of the methods M" that may be reached by a dynamic dispatch:

adding any type that occurs in the <u>set of</u> types  $S_M$  and that is a subtype of the type of a parameter of the methods M'' to athe set of types  $S_{M''}$ ;

adding any type that occurs in the set of types  $S_{M''}$  and that is a subtype of the return type of the methods M'' to the set of types  $S_{M}$ .

5. (Currently Amended) The program storage device according to claim 4, further the programming instructions of:

for each field F read by the method M, adding any type that occurs in the set of types  $S_F$  to the set of types  $S_M$ ; and

for each field F with the <u>set of all types</u> T written by the method M, add<u>ing</u> any type that occurs in the <u>set of types</u>  $S_M$  and that is a subtype of the <u>set of all types</u> T to the <u>set of types</u>  $S_F$ .

6. (Currently Amended) The program storage device according to claim 19, further comprising the programming instructions of:

using the call graph <del>computed above</del>, as previously constructed, in a compiler as a basis for performing optimizations such as inlining.

7. (Currently Amended) The program storage device according to claim 19, further comprising the programming instructions of:

using the call graph <del>computed above, as previously constructed,</del> in a reporting tool to report call graph information to a user.

8. (Currently Amended) A <u>computer program product program storage device, readable</u> by a machine, tangibly embodying instructions to perform method steps for constructing a <u>set of types occurring within a method call graph as a representation of a program</u>, the method computer program product comprising:

a storage medium readable by a processing circuit and storing instructions for execution by the processing circuit for performing a method comprising:

selecting a program P for constructing a call graph representation thereof; wherein the program P contains zero or more fields  $F_F$  and one or more at least two methods  $M_M$ ;

wherein each method  $M_1$  in  $M_M$  has a single body B; wherein for each method  $M_2$  in  $M_M$ , the call graph representation includes a corresponding node;

wherein the call graph representation includes zero or more edges corresponding to connections between two or more of nodes;

identifying constructing for each method M in  $M_M$ , a set of zero or more types  $S_M$  of objects which occur therein;

identifyingconstructing for each field F in F<sub>F</sub>, a set of zero or more types S<sub>F</sub> of objects stored therein;

identifying one or more allocation sites inside the body B of each of method M;

determining a set of directly called-methods M' that are directly called withininside the body B of each method M and propagating types from the set of types  $S_{M'}$  to the set of types  $S_{M}$  and from of the set of types  $S_{M}$  to the set of types  $S_{M'}$ ; and

determining a set of virtually called methods M" that are virtually called withininside the body B of each method M and propagating types from the set of types  $S_{M"}$  to the set of types  $S_{M}$  and from of the set of types  $S_{M}$  to the set of types  $S_{M}$ ;

## determining a set of fields F that are

read in the body B of each method M, and propagating types from the set of types  $S_F$  to the set of types  $S_M$ ; and

written in the body B of each method M, and propagating types from the set of types  $S_M$  to the set of types  $S_F$ .

9. (Currently Amended) The <u>computer program storage device product</u> according to claim 8, further comprising:

determining the set of all types T that are allocated in the body of method M, and adding each element of the set of all types T to the set of types  $S_M$  for each allocation of type T that occurs in the method M.

10. (Currently Amended) A method for constructing a <u>set of types occurring within a</u> method call graph as a representation of a program, the method comprising:

selecting a program P for constructing a call graph representation thereof; wherein the program P contains zero or more fields  $F_F$  and one or more at least two methods  $M_M$ ;

wherein each method M<sub>1</sub> in M<sub>M</sub> has a single body B;

- wherein for each method M<sub>2</sub> in M<sub>M</sub>, the call graph representation includes a corresponding node;
- wherein the call graph representation includes zero or more edges corresponding to connections between two or more of nodes;

identifying constructing for each method M in  $M_M$ , a set of zero or more types  $S_M$  of objects which occur therein;

identifying constructing for each field F in F<sub>F</sub>, a set of zero or more-types S<sub>F</sub> of objects stored therein;

identifying one or more allocation sites inside the body B of each of method M; determining a set of directly called methods M' that are directly called within inside the body B of each method M and propagating types from a set of types  $S_{M'}$  to the set of types  $S_{M}$  and from of the set of types  $S_{M}$  to the set of types  $S_{M'}$ ; and

determining a set of virtually called methods M" that are virtually called within inside the body B of each method M and propagating types from a set of types  $S_{M"}$  to the set of types  $S_{M}$  and from of the set of types  $S_{M}$  to the set of types  $S_{M"}$ ;

determining a set of fields F that are

read in the body B of each method M, and propagating types from the set of types  $S_F$  to the set of types  $S_M$ ; and

written in the body B of each method M, and propagating types from the set of types  $S_M$  to the set of types  $S_F$ .

11. (Currently Amended) The method according to claim 10, further comprising:

determining the set of all types T that are allocated in the body of method M, and adding each element of the set of all types T to the set of types S<sub>M</sub> for each allocation of

## type T that occurs in the method M.

12. (Currently Amended) The method according to claim 11, further comprising:

for each direct call to the <u>a</u> methods M' in the body B of the method M performing the steps of:

adding any type that occurs in the <u>set of</u> types  $S_M$  and that is a subtype of a type of a parameter of the methods M' to <u>the set of</u> types  $S_{M'}$ ; and adding any type that occurs in the <u>set of</u> types  $S_{M'}$  and that is a subtype of a return type of the methods M' to the <u>set of</u> types  $S_M$ .

13. (Currently Amended) The method according to claim 12, further comprising: for each virtual call to the methods M' in the body B of the method M:

using the <u>set of types</u>  $S_M$ , determine each of the methods M" that may be reached by a dynamic dispatch:

adding any type that occurs in the <u>set of</u> types  $S_M$  and that is a subtype of a type of a parameter of the methods M'' to athe set of types  $S_{M''}$ ;

adding any type that occurs in the set of types  $S_{M''}$  and that is a subtype of the return type of the methods M'' to the set of types  $S_{M}$ .

14. (Currently Amended) The method according to claim 13, further comprising: for each field F read by the method M, add any type that occurs in the set of types  $S_F$  to the set of types  $S_M$ ; and

for each field F with the <u>set of all</u> types T written by the method M, add any type that occurs in the <u>set of</u> types  $S_M$  and that is a subtype of the <u>with set of all types</u> T to the <u>set of</u> types  $S_F$ .

15. (Currently Amended) The method according to claim 10, further comprising the step of:

using the call graph <del>computed above</del>, <u>as previously constructed</u>, in a compiler as a basis for performing optimizations such as inlining.

16. (Currently Amended) The method according to claim 420, further comprising the step of:

using the call graph <del>computed above</del>, as previously constructed, in a reporting tool to report call graph information to a user.

- 17. (Cancelled)
- 18. (Cancelled)
- 19. (New) The program storage device according to claim 1, further comprising: constructing a call graph representation;

wherein for each method  $M_2$  in  $M_M$ , the call graph representation includes a corresponding node; and

wherein the call graph representation includes zero or more edges corresponding to connections between two or more of nodes.

20. (New) The method according to claim 10, further comprising: constructing a call graph representation;

wherein for each method  $M_2$  in  $M_M$ , the call graph representation includes a corresponding node; and

wherein the call graph representation includes zero or more edges corresponding to connections between two or more of nodes.